



Accelerating Polynomial Multiplication for Homomorphic Encryption on GPUs





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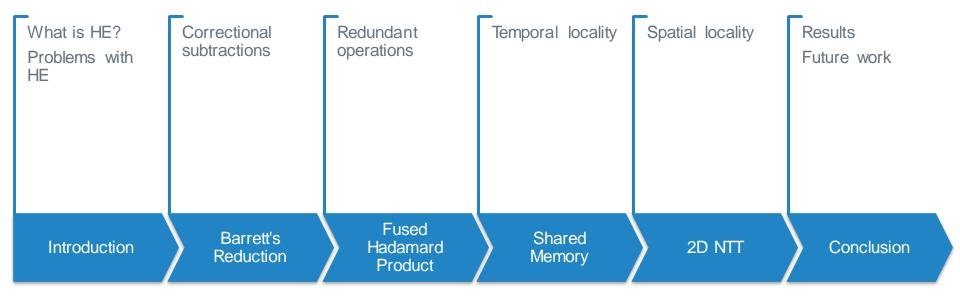






Outline

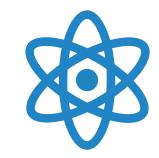




What is Homomorphic Encryption?



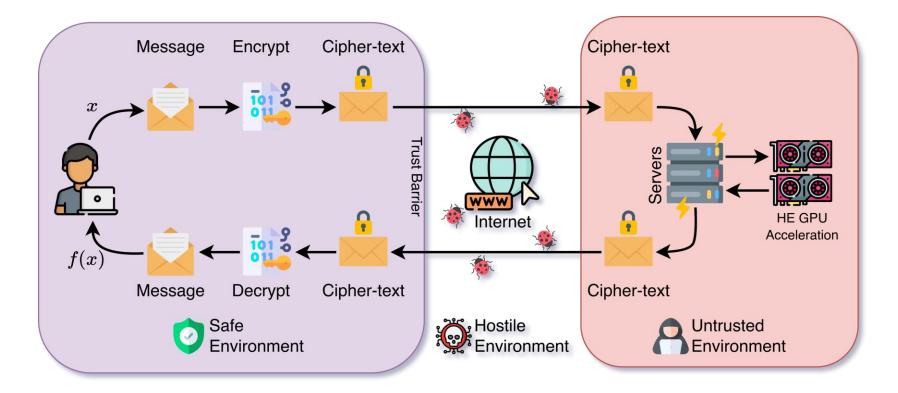




HE is a type of encryption that allows computation to be run on encrypted operands HE schemes are lattice-based, making them quantum resistant

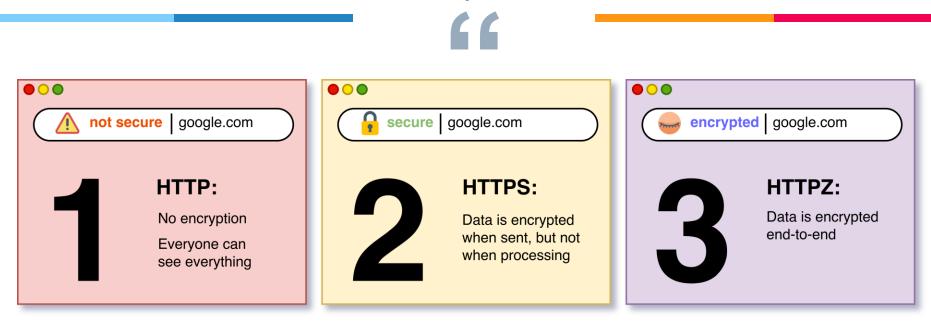
What is Homomorphic Encryption?







"HE could give rise to a new internet protocol, **HTTPZ**, that would standardize end-to-end encryption and replace HTTPS as the default protocol"

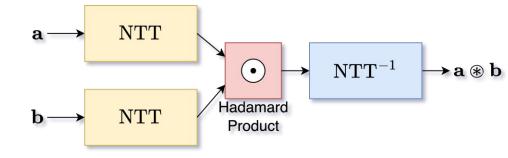


--- By R. Hindi, "People shouldn't care about privacy", ZAMA, Aug 2021

HE Bottlenecks



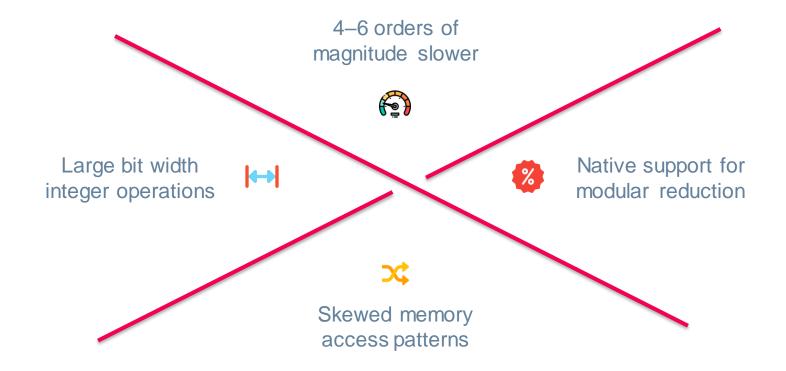
- Polynomial multiplication is the key bottleneck for lattice-based cryptography and HE
- Poly-mult is typically implemented with Number Theoretic Transform (NTT)
- NTT relies heavily on modular reduction operation



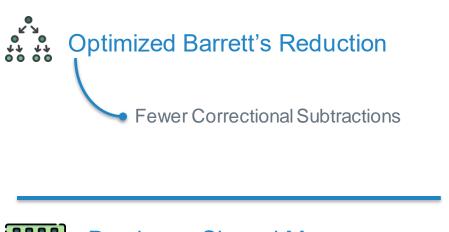
Polynomial Multiplication

Problems with HE





Contributions



Fused Hadamard Product

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Reduces Operational Complexity

Persistent Shared Memory

Increased Temporal Locality

Mixed Radix 2D NTT



Increased Spatial Locality

Optimized Barrett's Reduction



 $\operatorname{Remainder} = x \ \% \ q$

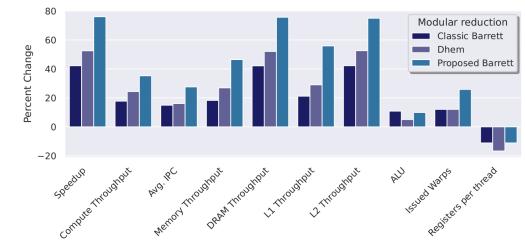
$$ext{Remainder} = x \; - \; \left(\left\lfloor rac{x}{q}
ight
floor imes q
ight)$$

- NTT is dominated by modulo operation
- Modulo computation involves expensive division operation
- **Barrett's reduction** replaces division with a set of bit-shift and multiplication operations

Algorithm 2 Classical Barrett reductionRequire: $m = \operatorname{len}(q) \leq \beta - 2, \ 0 \leq x < 2^{2m}, \ \mu = \lfloor \frac{2^{2m}}{q} \rfloor$ Ensure: $\operatorname{rem} = x \mod q$ 1: $c \leftarrow x \gg (m-1)$ 2: $\operatorname{quot} \leftarrow (c \times \mu) \gg (m+1)$ 3: $\operatorname{rem} \leftarrow x - \operatorname{quot} \times q$ 4: if $\operatorname{rem} \geq q$ then5: $\operatorname{rem} \leftarrow \operatorname{rem} - q$ 6: if $\operatorname{rem} \geq q$ then7: $\operatorname{rem} \leftarrow \operatorname{rem} - q$ 8: return rem

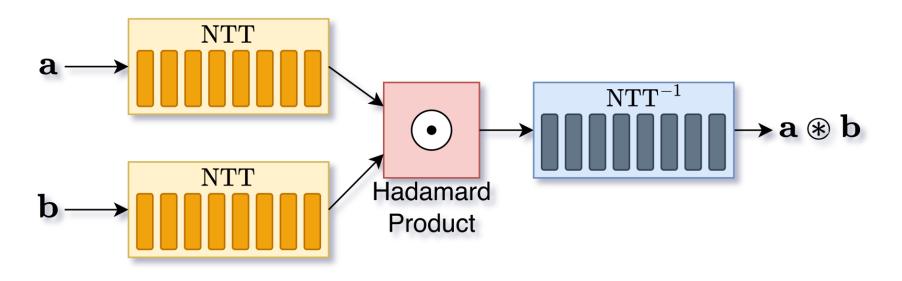
Optimized Barrett's Reduction

- Proposed Barrett's reduction reduces number of correctional subtractions
- Future Work: Natively supported modular reduction with hardware implementation

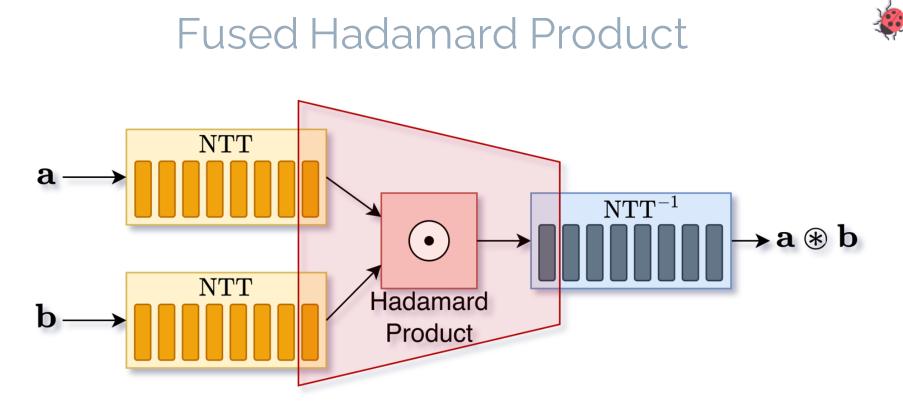






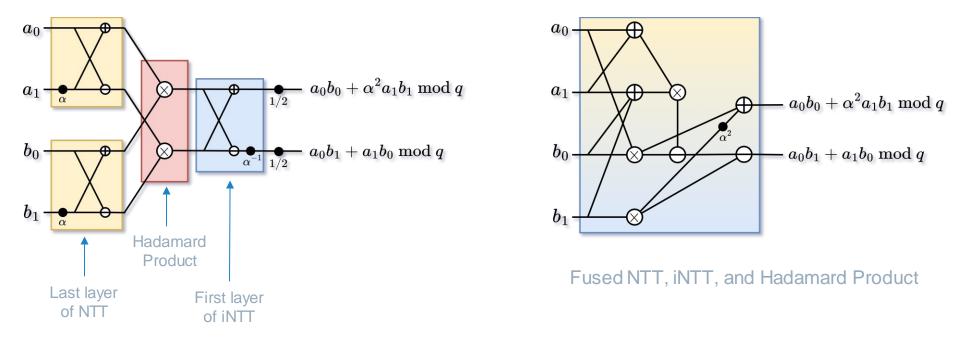


Polynomial Multiplication



Polynomial Multiplication

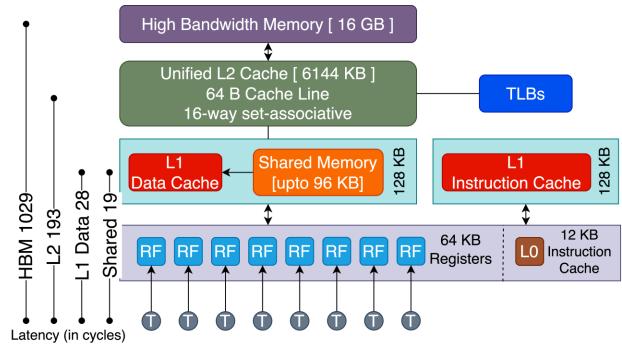
Fused Hadamard Product



Persistent Shared Memory



- NTT is a memory-bound kernel
- Each stage of NTT generates intermediate results for subsequent stages
- Intermediate results of each stage are cached on shared memory
- Removes redundant global memory accesses

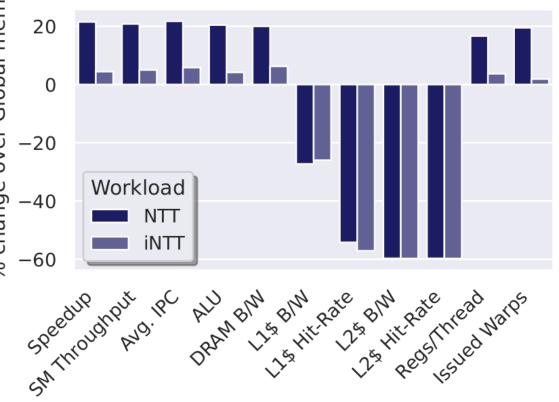


V100 GPU Memory Hierarchy

Persistent Shared Memory

% change over Global mem Shared memory use provides 25% -20speedup over global memory L1 and L2 cache memory pressure drops significantly -40Only works for NTT sizes that fit in shared memory size (N<2¹¹)

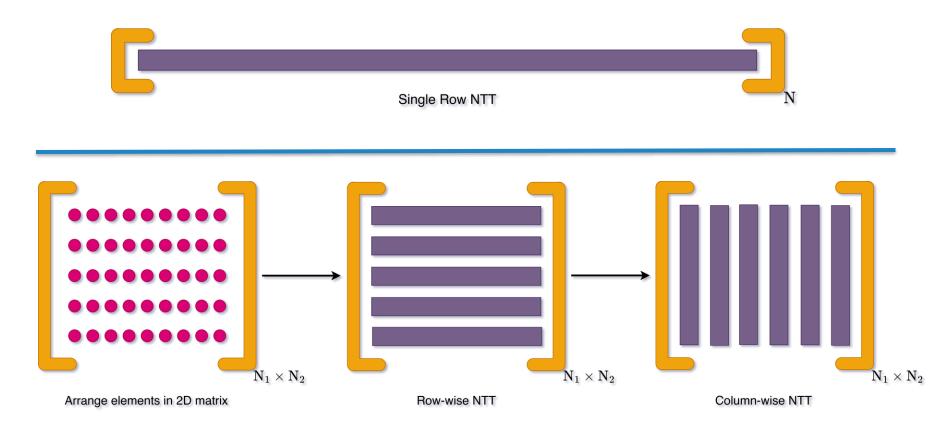
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Mixed Radix 2D NTT

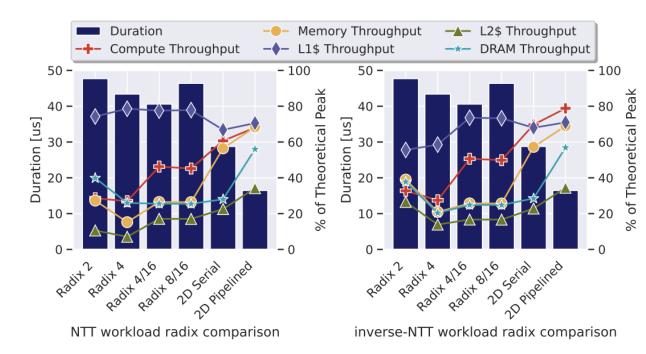




Mixed Radix 2D NTT



- 2D NTT increases compute costs
- Reduces memory pressure
- Preserves spatial locality
- Row-wise and Column-wise NTT can be pipelined



Concluding Remarks



HE is a **popular** memory-intensive workload with high computational demands Explored key bottlenecks in HE



- Algorithmic improvements
- Low-level kernel improvements

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- Presented four optimizations
 - Achieved speedup
 - CPU: 123.13x
 - GPU: 2.37x

Future Work / Architectural Feature Requests



- Larger integer bit width support for GPU
- Native modular reduction support





Thank you!

Any questions?

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